

## MATH 245 F24, Exam 2 Solutions

1. Carefully define the following terms: Proof by Shifted Induction, recurrence.

Let  $s \in \mathbb{Z}$ . To prove “ $\forall x \in \mathbb{Z}$  with  $x \geq s$ ,  $P(x)$ ” by shifted induction, we must (a) prove  $P(s)$ ; and (b) prove  $\forall x \in \mathbb{Z}$  with  $x \geq s$ ,  $P(x) \rightarrow P(x+1)$ . A recurrence is a **sequence** in which all but finitely many terms are defined out of the previous terms.

2. Carefully state The Master Theorem.

Let  $a \in \mathbb{N}$ ,  $k \in \mathbb{R}$ , and  $b \in \mathbb{R}$  with  $b > 1$  (three constants). Let  $c_n$  be a sequence with  $c_n = \Theta(n^k)$ . Set  $d = \log_b a$ . Consider now the sequence  $T_n = aT_{n/b} + c_n$ . The Master Theorem now has three cases:

If  $k < d$ , then  $T_n = \Theta(n^d)$ . If  $k = d$ , then  $T_n = \Theta(n^d \log n)$ . If  $k > d$  then  $T_n = \Theta(n^k)$ .

3. Let  $a_n = n^3 + 4n + \sin(n)$ . Prove or disprove that  $a_n = \Theta(n^3)$ .

The statement is true, and its proof has two halves.

$a_n = O(n^3)$ : Choose  $M = 6$  and  $n_0 = 1$ . Let  $n \geq n_0$ . Now,  $n^3 \geq n$  so  $4n^3 \geq 4n$ . Also,  $n^3 \geq n \geq 1$ . Now,  $|a_n| \leq |n^3| + |4n| + |\sin n| \leq n^3 + 4n + 1 \leq n^3 + 4n^3 + n^3 = 6n^3 = 6|b_n|$ .

$a_n = \Omega(n^3)$ : Choose  $M = 1$  and  $n_0 = 1$ . Let  $n \geq n_0$ . Now,  $\sin n \geq -1$ , so  $n + \sin n \geq n - 1 \geq 0$ , and hence  $n^3 + 3n + (n + \sin n) \geq 0$ . Hence,  $1|a_n| = |n^3 + 4n + \sin n| = n^3 + 4n + \sin n = n^3 + 3n + (n + \sin n) \geq n^3 + 0 + 0 = |b_n|$ .

COMMON ALGEBRA ERRORS: While it is true that  $|a + b| \leq |a| + |b|$  (this is called the triangle inequality), it is NOT true that  $|a + b| = |a| + |b|$ . For example, take  $a = 3, b = -3$ ; we have  $|a + b| = 0$  while  $|a| + |b| = 6$ . It is also NOT true that  $|\sin n| = \sin n$ . For example, take  $n = -\pi/2$ ; we have  $|\sin n| = 1$  while  $\sin n = -1$ . Some of you got these errors out of your system when you were solving exercise 7.15, but some of you did not.

4. Let  $n \in \mathbb{Z}$ . Prove that  $\frac{n^2 - n - 2}{2} \in \mathbb{Z}$ .

Note that  $n^2 - n - 2 = (n - 2)(n + 1)$ . We apply the Division Algorithm Theorem to  $n, 2$ , getting  $q, r \in \mathbb{Z}$  with  $n = 2q + r$  and  $0 \leq r < 2$ .

(or we can cite Cor. 1.8 to know that  $n$  is even or odd, then use the definitions of even, odd).

If  $n = 2q + 0$ , then  $\frac{n^2 - n - 2}{2} = \frac{(2q-2)(n+1)}{2} = (q-1)(n+1) \in \mathbb{Z}$ .

If instead  $n = 2q + 1$ , then  $\frac{n^2 - n - 2}{2} = \frac{(n-2)(2q+1+1)}{2} = (n-2)(q+1) \in \mathbb{Z}$ .

In both cases,  $\frac{n^2 - n - 2}{2} \in \mathbb{Z}$ .

5. Prove that for all  $n \in \mathbb{N}$ , we must have  $4^n \geq 3^n + 1$ .

Proof by (vanilla) induction. Base case,  $n = 1$ ,  $4^n = 4 \geq 3 + 1 = 3^n + 1$ .

Inductive case, let  $n \in \mathbb{N}$  and assume that  $4^n \geq 3^n + 1$ . Multiply both sides by 4 to get  $4^{n+1} = 4 \cdot 4^n \geq 4 \cdot 3^n + 4 \geq 4 \cdot 3^n + 1 \geq 3 \cdot 3^n + 1 = 3^{n+1} + 1$ . Hence  $4^{n+1} \geq 3^{n+1} + 1$ .

ALTERNATE ALGEBRA: Multiply both sides by 3 to get  $3 \cdot 4^n \geq 3 \cdot 3^n + 3$ . Now,  $4^{n+1} = 4 \cdot 4^n > 3 \cdot 4^n$  and also  $3 \cdot 3^n + 3 > 3 \cdot 3^n + 1 = 3^{n+1} + 1$ . Combine everything to get  $4^{n+1} \geq 3^{n+1} + 1$ .

6. Prove that the Fibonacci numbers satisfy, for all  $n \in \mathbb{N}_0$ , that  $F_n F_{n+1} = \sum_{i=0}^n F_i^2$ .

Proof by (shifted) induction. Base case  $n = 0$ :  $F_0 F_1 = 0 \cdot 1 = 0$ , while  $\sum_{i=0}^0 F_i^2 = F_0^2 = 0^2 = 0$ . Inductive case, let  $n \in \mathbb{N}_0$  and assume that  $F_n F_{n+1} = \sum_{i=0}^n F_i^2$ . Add  $F_{n+1}^2$  to both sides, getting  $F_{n+1}^2 + F_n F_{n+1} = \sum_{i=0}^{n+1} F_i^2$ . Now,  $F_{n+1}^2 + F_n F_{n+1} = F_{n+1}(F_{n+1} + F_n) = F_{n+1} F_{n+2}$ . Combining, we get  $F_{n+1} F_{n+2} = \sum_{i=0}^{n+1} F_i^2$ .

7. Solve the recurrence with initial conditions  $a_0 = 3, a_1 = -1$ , and relation  $a_n = a_{n-1} + 6a_{n-2}$  (for  $n \geq 2$ ).

The characteristic polynomial is  $x^2 - x - 6 = (x - 3)(x + 2)$ , whose roots are 3, -2. Hence the general solution is  $a_n = A3^n + B(-2)^n$ . We now apply the initial conditions  $3 = a_0 = A3^0 + B(-2)^0 = A + B$  and  $-1 = a_1 = A3^1 + B(-2)^1 = 3A - 2B$ . We now solve the system  $\{3 = A + B, -1 = 3A - 2B\}$  to get  $A = 1, B = 2$ . Hence, the specific solution is  $a_n = 3^n + 2(-2)^n$ .

8. Prove or disprove:  $\forall x \in \mathbb{Z}, \exists y \in \mathbb{Z}, x = y^2 + y$ .

The statement is false; a counterexample consists of an  $x$  and two different  $y$ 's. Many choices are possible:  $(x = 0, y_1 = 0, y_2 = -1), (x = 2, y_1 = 1, y_2 = -2), (x = 6, y_1 = 2, y_2 = -3)$ , etc. In all cases, a complete solution must show that  $x = y_1^2 + y_1$  and  $x = y_2^2 + y_2$ .

One full solution: Take  $x = 2, y_1 = 1, y_2 = -2$ . We have  $y_1^2 + y_1 = 1^2 + 1 = 2 = x$  and also  $y_2^2 + y_2 = (-2)^2 + (-2) = 4 - 2 = 2 = x$ . Since  $y_1 \neq y_2$ , the statement is false.

9. Let  $x, y \in \mathbb{R}$ . Without using any theorems from Chapter 5, prove that: if  $x \leq y$  then  $\lceil x \rceil \leq \lceil y \rceil$ .

Direct proof: suppose that  $x \leq y$ . We apply the definition of ceiling twice:  $\lceil x \rceil - 1 < x \leq \lceil x \rceil$  and  $\lceil y \rceil - 1 < y \leq \lceil y \rceil$ . We combine the first inequality with  $x$  and the second inequality with  $y$ , to get the chain  $\lceil x \rceil - 1 < x \leq y \leq \lceil y \rceil$ , and hence  $\lceil x \rceil - 1 < \lceil y \rceil$ . Since  $\lceil x \rceil - 1$  and  $\lceil y \rceil$  are both integers, we can apply Theorem 1.12(a) from the book (you don't need to know the number, just that it's a theorem and it's not in chapter 5), getting  $\lceil x \rceil - 1 \leq \lceil y \rceil - 1$ . Adding 1 to both sides we get the desired  $\lceil x \rceil \leq \lceil y \rceil$ .

10. Let  $a, b \in \mathbb{Z}$  with  $b \geq 1$ . Use maximum element induction to prove that there are  $q, r \in \mathbb{Z}$  satisfying  $a = 2bq + r$  and  $-b < r \leq b$ .

We set  $S = \{m \in \mathbb{Z} : m < \frac{a+b}{2b}\}$ . The real number  $\frac{a+b}{2b}$  is an upper bound for  $S$ , and  $S$  is nonempty because it is a halfline. By maximum element induction, there is some maximum to  $S$ , i.e. some integer  $q$  such that  $q < \frac{a+b}{2b}$  but  $q + 1 \geq \frac{a+b}{2b}$ . Set  $r = a - 2bq$ . We have  $2bq < a + b$  and hence  $-b < a - 2bq = r$ . We also have  $2bq + 2b \geq a + b$  so  $b \geq a - 2bq = r$ . Combining, we get  $-b < r \leq b$ , as desired. Since  $r = a - 2bq$ , we also have  $a = 2bq + r$ .